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Conditional independences in Gaussian vectors and rings of polynomials

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Abstract. Inference among the conditional independences in nondegenerate Gaussian vectors is studied by algebraic techniques. A general method to prove implications involving the conditional independences is presented. The method relies on computation of a Groebner basis. Examples of the implications are discussed.

1. Introduction

Conditional independence (CI-) constraints have become crucial ingredient in building highly structured probabilistic models of artificial intelligence. They provide a language communicating qualitative features and induce decompositions of probability distributions amenable to practical computations. Though the major part of research has concerned CI-constraints related to graphs, cf. [11, 2], an effort is devoted also to the general problem of discrete structures describing simultaneous occurance of independences within a random vector, for references see [16, 7]. Challenging problems involving CI-constraints occur also in statistics [5], information theory [17], secret-sharing schemes of cryptography [1], representations of matroid theory [9] and instances of the probabilistic method in combinatorics [14].

For a random vector $\xi = (\xi_i)_{i\in N}$ indexed by a finite set N, the stochastic conditional independences $\xi_I \perp\!\!\!\perp \xi_J|\xi_K$ among its subvectors $\xi_I = (\xi_i)_{i\in I}$, ξ_J, ξ_K, where $I, J, K \subseteq N$ are disjoint, can be encoded by means of the discrete structure

$$[\xi] = \{(ij|K) \in \mathcal{R};\ \xi_i \perp\!\!\!\perp \xi_j|\xi_K\},$$

see [6, Lemma 3]. Here, ij denotes a two-element set $\{i, j\}$ and $\mathcal{R}$ is the family of all ordered pairs $(ij|K)$ with $ij, K \subseteq N$ disjoint. Manipulations with the CI-constraints can

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be facilitated when understanding which subset $\mathcal{K}$ of $\mathcal{R}$ can be represented as $\mathcal{K} = [\xi]$ by means of some ξ. This difficult problem remains open even in a discrete setting with N having more than four elements, see [8, 7]. The same problem for Gaussian vectors has not been yet investigated in depth. This note aims at first steps in this direction, restricting attention to nondegenerate vectors. The degenerate case seems to differ only by additional technicalities.

Let $A = (a_{i,j})_{i,j=1}^{n}$ be the covariance matrix of a Gaussian vector $\xi = (\xi_1, \ldots, \xi_n)$, $n \geqslant 1$. When the vector is nondegenerate the matrix A is positive definite, in symbols $A > 0$. The variable ξ_1 is conditionally independent of ξ_2 given $(\xi_3, \ldots, \xi_n)$ if and only if the $(1,2)$-element of the inverse A^{-1} of A vanishes. Denoting submatrices of A by $A_{I,J} = (a_{i,j})_{i \in I, j \in J}$ where $I, J \subseteq N = \{1, \ldots, n\}$, this is equivalent to vanishing of the determinant $|A_{N\setminus\{1\},N\setminus\{2\}}|$, see for example [12, p. 11]. More generally, a conditional independence $\xi_i \perp\!\!\!\perp \xi_j | \xi_K$ for $(ij|K) \in \mathcal{R}$ takes place if and only if $|A_{jK,iK}|$, equal to $|A_{iK,jK}|$, vanishes; here iK abbreviates $\{i\} \cup K$ and similarly with jK. Since the independences are determined by the covariance matrix A of ξ one can work directly with positive definite matrices and use the notation

$$\langle\!\langle A \rangle\!\rangle = \{(ij|K) \in \mathcal{R}\,;\ |A_{iK,jK}| = 0\}$$

on the place of $[\xi]$. The main problem is now to recognize whether a given subset $\mathcal{K}$ of $\mathcal{R}$ can be written as $\mathcal{K} = \langle\!\langle A \rangle\!\rangle$ with some $A > 0$. An equivalent problem is to find for which families $\mathcal{L}, \mathcal{M} \subseteq \mathcal{R}$ the implication

$$\mathcal{L} \subseteq \langle\!\langle A \rangle\!\rangle \quad \Rightarrow \quad \mathcal{M} \cap \langle\!\langle A \rangle\!\rangle \neq \emptyset \tag{1}$$

holds for all $A > 0$. In fact, $\mathcal{K} = \langle\!\langle A \rangle\!\rangle$ with some $A > 0$ if and only if (1) with $\mathcal{L} = \mathcal{K}$ and $\mathcal{M} = \mathcal{R} \setminus \mathcal{K}$ is violated by some $A > 0$.

By Sylvester's criterion [12], the set of positive definite matrices is a semialgebraic subset of a Euclidean space, that is, a set defined by a finite number of polynomial inequalities. To decide whether for given $\mathcal{L}, \mathcal{M}$ the implication (1) quantified by 'for all $A > 0$' is true or false amounts to decide whether the semialgebraic set

$$\big\{A > 0\,;\ |A_{iK,jK}| = 0 \text{ once } (ij|K) \in \mathcal{L}, \text{ and } |A_{iK,jK}| \neq 0 \text{ once } (ij|K) \in \mathcal{M}\big\}$$

is empty or not. Algorithms of considerable complexity for manipulations with semialgebraic sets are available, see [10], and can recognize whether a given semialgebraic set is empty. Small instances of some statistical problems were attacked by such algorithms in [4], reporting successful treatment of a trivial example of the implication (1), see (3) below. Nevertheless, the general algorithms seem to provide no insight to structure of this special problem.

Since the construction $A \mapsto \langle\!\langle A \rangle\!\rangle$ makes sense even for a fairly general matrix, for example, for any symmetric matrix, the quantification $A > 0$ of the main problem can be replaced by the assumption that A belongs to another set of symmetric matrices. Modified problems obtained in this way approximate the main problem and are of theoretical interest *per se* because they enrich insight into abstract CI structures. A convenient replacement of $A > 0$ may require A to belong to an algebraic set (a set defined by polynomial equalities) or a constructible set of matrices (finite union of differences of algebraic sets). When matrices are admitted to have the complex numbers for elements

numerous methods of classical algebraic geometry, supported by computational procedures for Groebner bases, are at disposal.

For example, when $A > 0$ the principal minors $|A_{I,I}|$ are positive for $\emptyset \neq I \subseteq N$, and a straightforward relaxation of the quantification $A > 0$ is to require A to be in the set $\mathbb{A}$ of complex symmetric matrices having all proper principal minors nonzero. The discrete structures $\langle\!\langle A \rangle\!\rangle$ with $A \in \mathbb{A}$ keep reasonable properties and remain within the class of semigraphoids, which has been a framework for CI relations, see [11, 5, 7]. When $n = 3$ this relaxation works very well and provides the same class of structures $\langle\!\langle A \rangle\!\rangle$ as if $A > 0$, see Example 1. However, when $n \geqslant 4$ this is not the case, see Example 3. In general, an 'exact' relaxation, of purely theoretical interest, to a constructible set of matrices exists. In fact, disjoint union of the sets of matrices

$$\left\{A \in \mathbb{A};\ |A_{iK,jK}| = 0 \ \text{ for } \ (ij|K) \in \mathcal{K}, \ \text{ and } \ |A_{iK,jK}| \neq 0 \ \text{ for } \ (ij|K) \in \mathcal{R} \setminus \mathcal{K}\right\}$$

over all $\mathcal{K} \subseteq \mathcal{R}$ admitting the representation $\mathcal{K} = \langle\!\langle A \rangle\!\rangle$ for some $A > 0$ is a constructible set; obviously, (1) is valid for $A > 0$ if and only if it is valid for A in the union.

This contribution is organized as follows. In Section 2, examples of the implication (1) are presented for complex matrices. When specialized to positive definite matrices, they summarize basic properties of CI in nondegenerate Gaussian vectors. In Section 3, the implication (1) is analysed when A belongs to the set of matrices that violate each of given polynomial equalities (such sets, for example, $\mathbb{A}$ itself, are constructible when working with a finite number of equalities). A necessary and sufficient condition for validity of such an implication is presented in Proposition 1. The condition can be practically verified by computation of a single Groebner basis. Section 4 discusses possible uses of this result: all basic implications of Section 2 can be obtained as consequences of Proposition 1, and two additional implications for CI in Gaussian vectors are derived by a more sophisticated relaxation.

2. Basic properties of conditional independence

In this section, basic properties of CI in nondegenerate Gaussian vectors are derived as consequences of a single matrix identity. It is convenient to abbreviate $A_{I,I}$ by A_I and take $|A_\emptyset| = 1$.

Lemma 1. *For any complex matrix A, and $L \subseteq N$, $k \in N \setminus L$, and $i, j \in N \setminus kL$*

$$(-1)^{[i,j;k]}\,|A_{ikL,jkL}||A_L| = |A_{iL,jL}||A_{kL}| - |A_{iL,kL}||A_{kL,jL}| \tag{2}$$

where $[i, j; k] = 1$ if k is between i and j and $[i, j; k] = 0$ otherwise.

Proof. When $L = \emptyset$ this is an obvious formula for determinant expansion. When $L \neq \emptyset$ one can suppose that A_L is invertible, the opposite case following afterwards by continuity. First, let all elements of L be smaller than minimum of the set $ijk = \{i, j, k\}$. The Schur complement of A_L in $A_{ikL,jkL}$ is the matrix $B_{ik,jk} = A_{ik,jk} - A_{ik,L}(A_L)^{-1}A_{L,jk}$. Then $|A_{ikL,jkL}|$ factors into $|A_L||B_{ik,jk}|$ by [12, Theorem 3.1.1a]. Similarly, the four determinants of the right-hand side of (2) can be factored as $|A_{iL,jL}| = |A_L||B_{i,j}|$, etc. Then (2) follows from the expansion formula for $|B_{ik,jk}|$.

If L and i, j, k are in a general position then a unique permutation π of $ijk \cup L$ exists such that $\pi(l_1) < \pi(l_2)$ for $l_1 < l_2$ with l_1, l_2 in L or in ijk, and such that the elements of $\pi(L) = \{\pi(l);\ l \in L\}$ are smaller than the minimum of $\pi(i)\pi(j)\pi(k)$. Let C be the matrix obtained by applying π to the rows and columns of A simultaneously. By the previous argument

$$(-1)^{[\pi(i),\pi(j);\pi(k)]}\,|C_{\pi(ikL),\pi(jkL)}||C_{\pi(L)}| = |C_{\pi(iL),\pi(jL)}||C_{\pi(kL)}| - |C_{\pi(iL),\pi(kL)}||C_{\pi(kL),\pi(jL)}|\,.$$

Then $[\pi(i), \pi(j); \pi(k)] = [i, j; k]$, $|C_{\pi(L)}| = |A_L|$ and $|C_{\pi(kL)}| = |A_{kL}|$. Denoting by n_i the number of all $l \in L$ greater than i, one can see that $|C_{\pi(ikL),\pi(jkL)}|$ is the $(-1)^{n_i+n_j}$ multiple of $|A_{ikL,jkL}|$. Similar observations concern the expressions $|C_{\pi(iL),\pi(jL)}|$ and $|C_{\pi(iL),\pi(kL)}||C_{\pi(kL),\pi(jL)}|$. Hence, (2) obtains in full generality. □

Corollary 1. *For any symmetric complex matrix A*

(3) $\{(ij|kL), (ij|L)\} \subseteq \langle\!\langle A \rangle\!\rangle \;\Rightarrow\; (ik|L) \in \langle\!\langle A \rangle\!\rangle$ or $(jk|L) \in \langle\!\langle A \rangle\!\rangle$,

(4) $\{(ij|kL), (ik|L)\} \subseteq \langle\!\langle A \rangle\!\rangle \;\Rightarrow\; (ij|L) \in \langle\!\langle A \rangle\!\rangle$, *once* $|A_{kL}| \neq 0$,

(5) $\{(ij|kL), (ik|L)\} \subseteq \langle\!\langle A \rangle\!\rangle \;\Rightarrow\; (ik|jL) \in \langle\!\langle A \rangle\!\rangle$, *once* $|A_L||A_{kL}| \neq 0$,

(6) $\{(ij|kL), (ik|jL)\} \subseteq \langle\!\langle A \rangle\!\rangle \;\Rightarrow\; (ij|L) \in \langle\!\langle A \rangle\!\rangle$, *once* $|A_{jkL}||A_L| \neq 0$,

(7) $\{(ij|L), (ik|L)\} \subseteq \langle\!\langle A \rangle\!\rangle \;\Rightarrow\; (ij|kL) \in \langle\!\langle A \rangle\!\rangle$, *once* $|A_L| \neq 0$.

Proof. Rewriting (2) for a symmetric matrix A

$$(-1)^{[i,j;k]}\,|A_{ikL,jkL}||A_L| - |A_{iL,jL}||A_{kL}| = -|A_{iL,kL}||A_{jL,kL}|\,,$$

$|A_{ikL,jkL}| = 0$ and $|A_{iL,jL}| = 0$ imply $|A_{iL,kL}||A_{jL,kL}| = 0$. This is the implication (3).

Along the same guidelines the identity (2) allows to conclude $(ij|L) \in \langle\!\langle A \rangle\!\rangle$ in (4).

To prove (5), one switches j and k in Lemma 1

$$(-1)^{[i,k;j]}\,|A_{ijL,jkL}||A_L| = |A_{iL,kL}||A_{jL}| - |A_{iL,jL}||A_{jL,kL}|\,. \tag{8}$$

and combines this equality with (4).

When $|A_{ikL,jkL}| = 0$ and $|A_{ijL,jkL}| = 0$ the left-hand sides in (2) and (8) vanish. Solving the two homogeneous equations for $|A_{iL,jL}|$ and $|A_{iL,kL}|$, their determinant is equal to $|A_{jL}||A_{kL}| - |A_{jL,kL}||A_{kL,jL}|$. Using Lemma 1 with $i = j$, this determinant equals $|A_{jkL}||A_L|$ and is nonzero by assumption. Hence, $|A_{iL,jL}| = 0$ and (6) follows.

The implication (7) follows directly from (2). □

A subset $\mathcal{K}$ of $\mathcal{R}$ is called *semigraphoid* if it fulfills

$$\{(ij|kL), (ik|L)\} \subseteq \mathcal{K} \iff \{(ik|jL), (ij|L)\} \subseteq \mathcal{K}\,, \tag{9}$$

cf. (4) and (5), and *pseudographoid* if it fulfills the implication $\Rightarrow$ in

$$\{(ij|kL), (ik|jL)\} \subseteq \mathcal{K} \iff \{(ij|L), (ik|L)\} \subseteq \mathcal{K}\,, \tag{10}$$

cf. (6). A semigraphoid which is a pseudographoid is called *graphoid*. (For the pseudographoid axiom compare [11, (3.6e) and (3.6b)] and [7, p. 107]. For (3) see [11, (3.34f)].)

By Corollary 1, to ensure that $\langle\!\langle A \rangle\!\rangle$ satisfies (9) and (10) it suffices to assume that $|A_I|$ is nonzero for every nonempty proper subset of N, or equivalently, $A \in \mathbb{A}$. Since all positive

definite matrices belong to $\mathbb{A}$ the conditional independences in nondegenerate Gaussian vectors enjoy (9), (10) and (3) with $\mathcal{K}$ or $\langle\!\langle A \rangle\!\rangle$ replaced by $[\xi]$.

Example 1. For $n = 3$, the implications (9), (10) and (3) show that $\langle\!\langle A \rangle\!\rangle$ with $A \in \mathbb{A}$ equals one of the following families $\emptyset$, $\{(ij|\emptyset)\}$, $\{(ij|k)\}$, $\{(ij|\emptyset), (ik|\emptyset), (ij|k), (ik|j)\}$, and $\mathcal{R}$. All five situations are possible, taking even the positively definite matrices ($\epsilon \neq 0$ sufficiently small)

$$\begin{bmatrix} 1 & \epsilon & \epsilon^2 \\ \epsilon & 1 & \epsilon^4 \\ \epsilon^2 & \epsilon^4 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & \epsilon \\ 0 & 1 & \epsilon \\ \epsilon & \epsilon & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & \epsilon & \epsilon^2 \\ \epsilon & 1 & \epsilon^3 \\ \epsilon^2 & \epsilon^3 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & \epsilon \\ 0 & \epsilon & 1 \end{bmatrix}$$

and the identity matrix, respectively. Therefore, the class of structures $\langle\!\langle A \rangle\!\rangle$ with $A > 0$ coincides with the class of $\langle\!\langle A \rangle\!\rangle$ with $A \in \mathbb{A}$. This is no longer true when $n > 3$, see Example 3.

3. Implications and polynomial rings

In this section a necessary and sufficient condition for validity of (1) when A is in a certain constructible set of matrices is presented. The condition involves rings of polynomials, for more details see [3, 10].

When studying $\langle\!\langle A \rangle\!\rangle$ for a symmetric matrix A with nonzero diagonal elements there is no loss of generality in assuming that all diagonal elements are equal to 1. In fact, multiplying each i-th row and i-th column of A by a number b_i satisfying $b_i^2 = a_{i,i}^{-1}$ one obtains a matrix A' rendering $\langle\!\langle A \rangle\!\rangle = \langle\!\langle A' \rangle\!\rangle$; in addition, A' can be made positive definite or belongs to $\mathbb{A}$ if A was positive definite or belonged to $\mathbb{A}$, respectively. From now on all matrices are supposed to be symmetric and to have the ones on their diagonals.

Let x_{ij} be an indeterminate indexed by a two-element set ij, $1 \leqslant i < j \leqslant n$, and $\mathfrak{R}$ be the ring of polynomials in these indeterminates with coefficients in the field of complex numbers. By little abuse of language we will substitute complex matrices A to polynomials, having in mind that the indeterminates x_{ij} are replaced by the off-diagonal elements $a_{i,j} = a_{j,i}$. For $(ij|K) \in \mathcal{R}$ the formula expressing the determinant $|A_{iK,jK}|$ by elements of $A_{iK,jK}$ gives rise to a polynomial of $\mathfrak{R}$. Denoting the polynomial by $f_{(ij|K)}$, one has $f_{(ij|K)}(A) = |A_{iK,jK}|$. The notation $f_{(ij|K)}$ is extended to $f_{\mathcal{M}} = \prod_{(ij|K)\in\mathcal{M}} f_{(ij|K)}$, $\mathcal{M} \subseteq \mathcal{R}$. Similarly, the formula for $|A_I|$ gives rise to a polynomial denoted by f_I. For example, $f_{(12|3)} = x_{12} - x_{13}x_{23}$, $f_{(13|2)} = x_{12}x_{23} - x_{13}$ and $f_{12} = 1 - x_{12}^2$.

Let p_r, $1 \leqslant r \leqslant t$, be some polynomials in $\mathfrak{R}$ and $\mathbb{B}$ be the set of complex matrices A (symmetric, with ones on diagonals) such that $p_r(A) \neq 0$ for all $1 \leqslant r \leqslant t$. Where y_r are additional indeterminates, the polynomials $1 - y_r p_r$ belong to the ring $\mathfrak{T} = \mathfrak{R}[y_1, \ldots, y_t]$ of polynomials in both x_{ij}'s and y_r's. Obviously, a matrix A together with complex numbers $b_1, \ldots, b_t$ solve the equations $1 - y_r p_r = 0$ if and only if $A \in \mathbb{B}$ and $b_r = p_r(A)^{-1}$. Thus $\mathbb{B}$ is a coordinate projection of an algebraic set and is therefore constructible by [3, pp. 258–9]. Another indeterminate z and the ring $\mathfrak{T}[z]$ of polynomials with one more indeterminate come into play.

Proposition 1. *Let $\mathcal{L}, \mathcal{M}$ be subfamilies of $\mathcal{R}$, p_r be a polynomial in $\mathfrak{R}$, $1 \leqslant r \leqslant t$, and $\mathbb{B}$ be the corresponding set of complex matrices. The implication*

$$\mathcal{L} \subseteq \langle\!\langle A \rangle\!\rangle \quad \Rightarrow \quad \mathcal{M} \cap \langle\!\langle A \rangle\!\rangle \neq \emptyset \,, \qquad A \in \mathbb{B} \,, \tag{11}$$

is true if and only if the ring $\mathfrak{T}[z]$ is generated by the polynomials $f_{(ij|K)}$ for $(ij|K) \in \mathcal{L}$, $1 - y_r p_r$ for $1 \leqslant r \leqslant t$, and $1 - z f_{\mathcal{M}}$. This is equivalent to the incidence of $f_{\mathcal{M}}$ to the radical $\sqrt{I_{\mathcal{L}}}$ of the ideal $I_{\mathcal{L}}$ generated by $f_{(ij|K)}$ for $(ij|K) \in \mathcal{L}$, and $1 - y_r p_r$ for $1 \leqslant r \leqslant t$.

Proof. If $\mathfrak{T}[z]$ is generated as above then the constant polynomial 1 can be combined from the listed polynomials. Thus, polynomials $h_{(ij|K)}$, h_r, and h in $\mathfrak{T}[z]$ exist such that

$$\sum_{(ij|K) \in \mathcal{L}} h_{(ij|K)} f_{(ij|K)} + \sum_{r=1}^{t} h_r [1 - y_r p_r] + h[1 - z f_{\mathcal{M}}] = 1 \,.$$

Substituting $A \in \mathbb{B}$ for x_{ij}'s and the nonzero numbers $p_r(A)^{-1}$ for y_r's the sum over r equals zero. Then

$$\sum_{(ij|K) \in \mathcal{L}} h^*_{(ij|K)} |A_{iK,jK}| + h^* \Big[1 - z \prod_{(ij|K) \in \mathcal{M}} |A_{iK,jK}| \Big] = 1$$

where $h^*_{(ij|K)}$ and h^* are polynomials involving only the indeterminate z. If $\langle\!\langle A \rangle\!\rangle \subseteq \mathcal{L}$ the sum over $\mathcal{L}$ vanishes. One concludes that the product must be equal to zero, that is, $\mathcal{M} \cap \langle\!\langle A \rangle\!\rangle \neq \emptyset$ and (11) holds.

Suppose (11) is true. When the polynomials $f_{(ij|K)}$ for $(ij|K)$ in $\mathcal{L}$, and $1 - y_r p_r$ for $1 \leqslant r \leqslant t$ vanish after substitution of complex numbers a_{ij} for x_{ij} and b_r for y_r then the symmetric matrix A with the off-diagonal elements $a_{i,j} = a_{ij}$ (and ones on the diagonal) belongs to $\mathbb{B}$. Now, $0 = f_{(ij|K)}(A) = |A_{iK,jK}|$ for $(ij|K)$ in $\mathcal{L}$ and thus $\langle\!\langle A \rangle\!\rangle \subseteq \mathcal{L}$. Due to (11), one has $(ij|K) \in \langle\!\langle A \rangle\!\rangle$ for some $(ij|K) \in \mathcal{M}$. Hence, $\prod_{(ij|K) \in \mathcal{M}} |A_{iK,jK}|$ equals zero and the polynomial $f_{\mathcal{M}}$ vanishes by the same substitution, $f_{\mathcal{M}}(A) = 0$. Hilbert's Nullstellensatz [3, p. 170] implies $f_{\mathcal{M}} \in \sqrt{I_{\mathcal{L}}}$.

If $f_{\mathcal{M}}$ belongs to the radical then polynomials $h_{(ij|K)}$ and h_r from $\mathfrak{T}$ exist such that

$$f^m_{\mathcal{M}} = \sum_{(ij|K) \in \mathcal{L}} h_{(ij|K)} f_{(ij|K)} + \sum_{r=1}^{t} h_r [1 - y_r p_r]$$

for an integer $m \geqslant 1$. Then $1 = z^m f^m_{\mathcal{M}} + [1 - z^m f^m_{\mathcal{M}}] = z^m f^m_{\mathcal{M}} + h^*[1 - z f_{\mathcal{M}}]$ where $h^* \in \mathfrak{T}[z]$. As $z^m f^m_{\mathcal{M}}$ is in $I_{\mathcal{L}}$ the ring $\mathfrak{T}[z]$ can be generated by $I_{\mathcal{L}}$ and $1 - z f_{\mathcal{M}}$, cf. [3, p. 176]. □

4. Discussion and examples

When investigating an instance of (1) for $A > 0$ Proposition 1 can be useful. Indeed, if each polynomial p_r is chosen so that $p_r(A) \neq 0$ for every $A > 0$ with ones on its diagonal Proposition 1 provides a sufficient condition for validity of the instance.

For example, since (2) implies

$$-(-1)^{[i,j;k]} z f_L f_{(ij|kL)} + z f_{kL} f_{(ij|L)} + \big[1 - z f_{(ik|L)} f_{(jk|L)}\big] = 1$$

Proposition 1 with $\mathcal{L} = \{(ij|kL), (ij|L)\}$, $\mathcal{M} = \{(ik|L), (jk|L)\}$ and $t = 0$ proves validity of (3) for matrices with nonzero diagonal elements. It is not difficult to see (and we do not give details) that all implications of Corollary 1 quantified by $A \in \mathbb{A}$ can be obtained

in this way from Proposition 1 with $t = 1, 2$ and f_I's in the role of p_r's. Hence, all basic properties of CI in nondegenerate Gaussian vectors can be viewed as consequences of Proposition 1.

When $n \geqslant 4$ a more sophisticated relaxation works. Given a matrix A and a permutation π of the set N, the matrix $A^\pi = (a_{\pi(i),\pi(j)})_{i,j=1}^n$ is obtained by permuting rows and columns of A simultaneously. Obviously, $A > 0$ if and only if $A^\pi > 0$. By [12, Theorem 36.2.1] the Hadamard product $A \circ A^\pi = (a_{i,j}\, a_{\pi(i),\pi(j)})_{i,j=1}^n$ of $A > 0$ and A^π is positive definite. Hence, the principal minors of Hadamard products $A^{\pi_1} \circ \ldots \circ A^{\pi_k}$ are positive for $A > 0$, $k \geqslant 1$ and arbitrary permutations $\pi_1, \ldots, \pi_k$ of N. As demonstrated below these minors are suitable polynomials to consider in Proposition 1 in the role of p_r's.

Example 2. To see that the implication

$$\{(12|3), (13|4), (14|2)\} \subseteq \langle\!\langle A \rangle\!\rangle \quad \Rightarrow \quad (12|\emptyset) \in \langle\!\langle A \rangle\!\rangle \tag{12}$$

for $A > 0$ is a consequence of Proposition 1, observe that $1 - a_{23}^2 a_{24}^2 a_{34}^2$ is a principal minor of Hadamard product of three matrices $A \circ A^{\pi_1} \circ A^{\pi_2}$; note that the polynomials $f_{(12|3)}$, $f_{(13|4)}$, $f_{(14|2)}$, $1 - y_1(1 - x_{23}^2 x_{24}^2 x_{34}^2)$ and $1 - z x_{12}$ generate $\mathfrak{T}[z]$, due to

$$1 = (1 - z x_{12}) + z x_{12}\big(1 - y_1(1 - x_{23}^2 x_{24}^2 x_{34}^2)\big) + y_1 z x_{12}\big[1 - x_{23}^2 x_{24}^2 x_{34}^2\big]$$

and

$$x_{12}\big[1 - x_{23} x_{24} x_{34}\big] = f_{(12|3)} + x_{23} f_{(13|4)} - x_{23} x_{34} f_{(14|2)} \ .$$

Analogously,

$$\{(12|3), (23|4), (34|1), (41|2)\} \subseteq \langle\!\langle A \rangle\!\rangle \quad \Rightarrow \quad (12|\emptyset) \in \langle\!\langle A \rangle\!\rangle \tag{13}$$

holds for $A > 0$, using the principal minor $1 - a_{13}^2 a_{24}^2$ of Hadamard product of two matrices and the identity

$$x_{12}\big[1 - x_{13}^2 x_{24}^2\big] = f_{(12|3)} + x_{13} f_{(23|4)} + x_{13} x_{24} f_{(34|1)} - x_{13}^2 x_{24} f_{(14|2)} \ .$$

Example 3. For ϵ small and nonzero the matrix A given by

$$\begin{bmatrix} 1 & \epsilon^2 & \epsilon^4 & \epsilon^3 \\ \epsilon^2 & 1 & \epsilon^{-2} & \epsilon \\ \epsilon^4 & \epsilon^{-2} & 1 & \epsilon \\ \epsilon^3 & \epsilon & \epsilon & 1 \end{bmatrix} \qquad \text{or} \qquad \begin{bmatrix} 1 & \epsilon^2 & \epsilon & \epsilon \\ \epsilon^2 & 1 & \epsilon & \epsilon^{-1} \\ \epsilon & \epsilon & 1 & \epsilon^2 \\ \epsilon & \epsilon^{-1} & \epsilon^2 & 1 \end{bmatrix}$$

renders $\langle\!\langle A \rangle\!\rangle$ equal to $\{(12|3), (13|4), (14|2)\}$ or $\{(12|3), (23|4), (34|1), (41|2)\}$, respectively. By Example 2, these two subsets of $\mathcal{R}$ are not of the form $\langle\!\langle A \rangle\!\rangle$ with $A > 0$ (the two matrices belong to $\mathbb{A}$ but are not positive definite.) Thus, to exclude vanishing of proper minors is not sufficient for algebraic study of CI in Gaussian vectors.

Due to the following assertion, it is not difficult to lift the implications (12) and (13) of Example 2 to a general form.

Lemma 2. *Let $\mathcal{L}, \mathcal{M} \subseteq \mathcal{R}$ and $L \subseteq N$ such that $L \cap (ij \cup K) = \emptyset$ for every $(ij|K)$ in $\mathcal{L} \cup \mathcal{M}$. If (1) is true for all $A > 0$ or $A \in \mathbb{A}$ then it remains true also when $\mathcal{L}$ and $\mathcal{M}$ are replaced by $\mathcal{L}_L = \{(ij|K \cup L);\ (ij|K) \in \mathcal{L}\}$ and $\mathcal{M}_L$, respectively.*

A proof is simple and therefore omitted.

Corollary 2. *For any positively definite matrix A, $L \subseteq N$ and $i,j,k,l \in N \setminus L$ different*

$$\{(ij|kL),(ik|lL),(il|jL)\} \subseteq \langle\!\langle A \rangle\!\rangle \quad \Rightarrow \quad (ij|L) \in \langle\!\langle A \rangle\!\rangle, \tag{14}$$

$$\{(ij|kL),(jk|lL),(kl|iL),(li|jL)\} \subseteq \langle\!\langle A \rangle\!\rangle \quad \Rightarrow \quad (ij|L) \in \langle\!\langle A \rangle\!\rangle\ . \tag{15}$$

The implication (15) goes back to [13, 15], and for (14) see [7].

From the computational point of view, a ring of polynomials is generated by a finite list of polynomials if and only if any Groebner basis of the ideal generated by the list contains a constant (nonzero) polynomial. Therefore, a verification whether $\mathfrak{T}[z]$ is generated as in Proposition 1 has the same complexity as a computation of a single Groebner basis.[†]

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[†]For example, in the software package Mathematica, the following input

$$\texttt{GroebnerBasis}\big[\{\mathtt{x}_{12}-\mathtt{x}_{13}\mathtt{x}_{23},\mathtt{x}_{13}-\mathtt{x}_{14}\mathtt{x}_{34},\mathtt{x}_{12}\mathtt{x}_{24}-\mathtt{x}_{14},1-\mathtt{y}_1[1-\mathtt{x}_{23}^2\mathtt{x}_{24}^2\mathtt{x}_{34}^2],1-\mathtt{z}\,\mathtt{x}_{12}\},$$
$$\{\mathtt{x}_{12},\mathtt{x}_{13},\mathtt{x}_{14},\mathtt{x}_{23},\mathtt{x}_{24},\mathtt{x}_{34},\mathtt{y}_1,\mathtt{z}\}\big]$$

is evaluated as the list $\{1\}$ consisting of the polynomial 1, cf. Example 2.