

III Operator Methods

① "Closure Properties"

Let (K, D) be a differential field and $C = \text{Const } K$. Until now we considered the problem of "finding" explicit solutions y of a given ODE

$$y^{(r)} + a_{r-1}y^{(r-1)} + \dots + a_0y = 0$$

with coefficients in K .

Now: use the equation itself as an implicit description of its solution(s), either together with a suitable number of initial values of its Taylor expansion to separate it from the other solutions, or simply as a "generic" element of the solution space.

(Compare: implicit definition of an algebraic number by its minimal polynomial and, optionally, a sufficiently accurate approximation)

Motivation: If F is a space of "functions" (more precisely, a K -VS with a function $\partial: F \rightarrow F$ satisfying $\partial(a \cdot f) = \partial(a) \cdot f + a \cdot \partial(f)$ for all $a \in K, f \in F$) and we let elements of $K[\partial]$ act on elements of F via

$$(a_0 + a_1 \partial + \dots + a_n \partial^n) \circ f := a_0 f + a_1 \partial(f) + \dots + a_n \partial^n(f)$$

then we have

$$\underbrace{(L_1 \cdot L_2)}_{\substack{\uparrow \\ \text{mul in} \\ K[\partial]}} \circ \underbrace{f}_{\substack{\uparrow \\ \text{appl}}} = L_1 \circ \underbrace{(L_2 \circ f)}_{\substack{\uparrow \\ \text{appl}}}$$

for all $L_1, L_2 \in K[\partial]$ and $f \in F$.

Def:

(1) $f \in F$ is called a solution of $L \in K[\partial]$ $\Leftrightarrow L \circ f = 0$

(2) For $L \in K[\partial]$, $V(L) := \{f \in F \mid L \circ f = 0\}$ is called the solution space of L

(3) For $f \in F$, $\text{ann}(f) := \{L \in K[\partial] \mid L \circ f = 0\}$
is called the annihilator of f .

(4) $f \in F$ is called D-finite

$$:\Leftrightarrow \exists L \in K[\partial] \setminus \{0\} : L \circ f = 0.$$

(\Leftrightarrow The K -Vs generated by $f, \partial f, \partial^2 f, \dots$
has finite dimension)

Note:

(1) If $L = L_1 L_2$ and f is a solution of L_2
then f is also a solution of L ,
because $L \circ f = (L_1 L_2) \circ f = L_1 \circ \underbrace{(L_2 \circ f)}_{=0} = 0$.

(2) If $L = L_1 + L_2$ and f is a common
solution of L_1 and L_2 , then f is
also a solution of L , because $0+0=0$.

It follows that $\text{ann}(f)$ is a left ideal
in $K[\partial]$.

For any two $L_1, L_2 \in K[x]$ with $L_2 \neq 0$ we can find unique $Q, R \in K[x]$ with

$$L_1 = QL_2 + R \quad \text{and} \quad \deg R < \deg L_2$$

by repeatedly subtracting suitable multiples $ad^i L_2$ from L_1 , like in the commutative case. We write $\text{quo}(L_1, L_2) := Q$, $\text{rem}(L_1, L_2) := R$.

If I is a left ideal in $K[x]$ and $L_1, L_2 \in I$, then also $L_1 - QL_2 = R \in I$. It follows that if h denotes a minimal order element of $I \setminus \{0\}$, then any other element L of I can be written $L = Q \cdot h$, because $L = QL + R$ with some $R \neq 0$ would yield a lower order element R of I in contradiction to the choice of h .

Furthermore, if h_1, h_2 are two minimal order elements of $I \setminus \{0\}$, then $h_1 = h_2$, for otherwise $h_1 - h_2$ would be a lower order element.

It follows that for every ideal $I \subseteq K[\partial]$ there exists a unique monic operator $L \in I$ with $I = \langle L \rangle$. For given $L_1, L_2 \in K[\partial]$, a $L \in K[\partial]$ with $\langle L_1, L_2 \rangle = \langle L \rangle$ can be computed by a straightforward adaptation of the Euclidean algorithm. We call $L := \text{gcd}(L_1, L_2)$ the greatest common right divisor of L_1 and L_2 .

Observe: If $L = \text{gcd}(L_1, L_2)$ then

$$V(L) = V(L_1) \cap V(L_2).$$

Conversely, if $L \in K[\partial]$ is the unique monic operator with $\langle L \rangle = \langle L_1 \rangle \cap \langle L_2 \rangle$, then L is called the least common left multiple of L_1 and L_2 , $L := \text{lclm}(L_1, L_2)$, and we have

$$V(L) = V(L_1) + V(L_2),$$

i.e. the solutions of L are linear combinations of solutions of L_1 or L_2 .

Note: the formula $\text{lcm}(p, q) = \frac{p \cdot q}{\text{gcd}(p, q)}$ does not carry over to $K[\partial]$. But we can compute $\text{lcm}(L_1, L_2)$ as follows:

- Make an ansatz $L = l_0 + l_1 \partial + \dots + l_n \partial^n$ with undetermined coeffs l_i .
- Compute $L \bmod L_1 = \square + \square \partial + \dots + \square \partial^{\deg L_1 - 1}$
 $L \bmod L_2 = \square + \square \partial + \dots + \square \partial^{\deg L_2 - 1}$
- The \square are certain K -linear combinations of the undetermined coefficients l_i . Forcing them to zero gives a linear system of equations over K with $n+1$ variables and $\deg L_1 + \deg L_2$ equations. Solve it.
- If there is no nontrivial solution, set $n = n+1$ and try again.

The smallest n for which a solution is found gives the lcm. At latest this happens when $n = \deg L_1 + \deg L_2$.

In particular: The sum of two D-finite functions is D-finite.

If $f_1 \in U(L_1)$, $f_2 \in U(L_2)$ then $\text{lclm}(L_1, L_2) \in \text{ann}(f_1 + f_2)$,

but it may not be the generator.

Product. Suppose now that (F, ∂) is even a differential ring and let $f_1, f_2 \in F$ be D-finite, say $L_1 \circ f_1 = 0$, $L_2 \circ f_2 = 0$ for some $L_1, L_2 \in K[\partial]$ with $d_1 = \deg L_1$, $d_2 = \deg L_2$.

The derivatives $f_1, \partial f_1, \partial^2 f_1, \dots$ generate the same K-US as $f_1, \partial f_1, \dots, \partial^{d_1-1} f_1$. Likewise for f_2 . Consequently, the K-US generated by all the products $(\partial^i f_1)(\partial^j f_2)$ ($i, j \in \mathbb{N}$)

$$\begin{array}{cccc} f_1 f_2 & f_1 f_2' & f_1 f_2'' & \dots \\ f_1' f_2 & f_1' f_2' & f_1' f_2'' & \dots \\ f_1'' f_2 & f_1'' f_2' & f_1'' f_2'' & \dots \\ \vdots & \vdots & \vdots & \dots \end{array}$$

is in fact generated by

$(\partial^i f_1)(\partial^j f_2)$ ($i=0 \dots d_1-1; j=0 \dots d_2-1$), so the dimension is at most $d_1 d_2$. Now observe that

$$\begin{aligned} f_1 f_2 \\ \partial(f_1 f_2) &= (\partial f_1) f_2 + f_1 (\partial f_2) \\ \partial^2(f_1 f_2) &= (\partial^2 f_1) f_2 + 2(\partial f_1)(\partial f_2) + f_1 (\partial^2 f_2) \\ &\vdots \end{aligned}$$

all belong to this vector space. Therefore, the first $d_1 d_2 + 1$ must be linearly dependent over K . The linear dependence corresponds to an operator

$$\square + \square \partial + \dots + \square \partial^{d_1 d_2}$$

which annihilates $f_1 f_2$.

In particular, the product of D -finite functions is D -finite.

An annihilating operator for ~~$f_1 f_2$~~ $f_1 f_2$ can be computed from annihilating operators L_1, L_2 of f_1, f_2 by following

the steps of this argument.

The (unique) monic least order operator which can be found in this way is called the symmetric product of L_1 and L_2 . It does not depend on the choice $f_1 \in U(L_1)$, $f_2 \in U(L_2)$ but only on L_1 and L_2 .

If L is the symmetric product of L_1, L_2 then

~~$\forall f_1, f_2$~~

$$\forall f_1 \in U(L_1) \quad \forall f_2 \in U(L_2):$$

$$L \in \text{ann}(f_1 f_2).$$

But L need not be the generator of these annihilators,